

# P2Cast: Peer-to-Peer Patching Scheme for VoD Service

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presented by Kyoungwon Suh

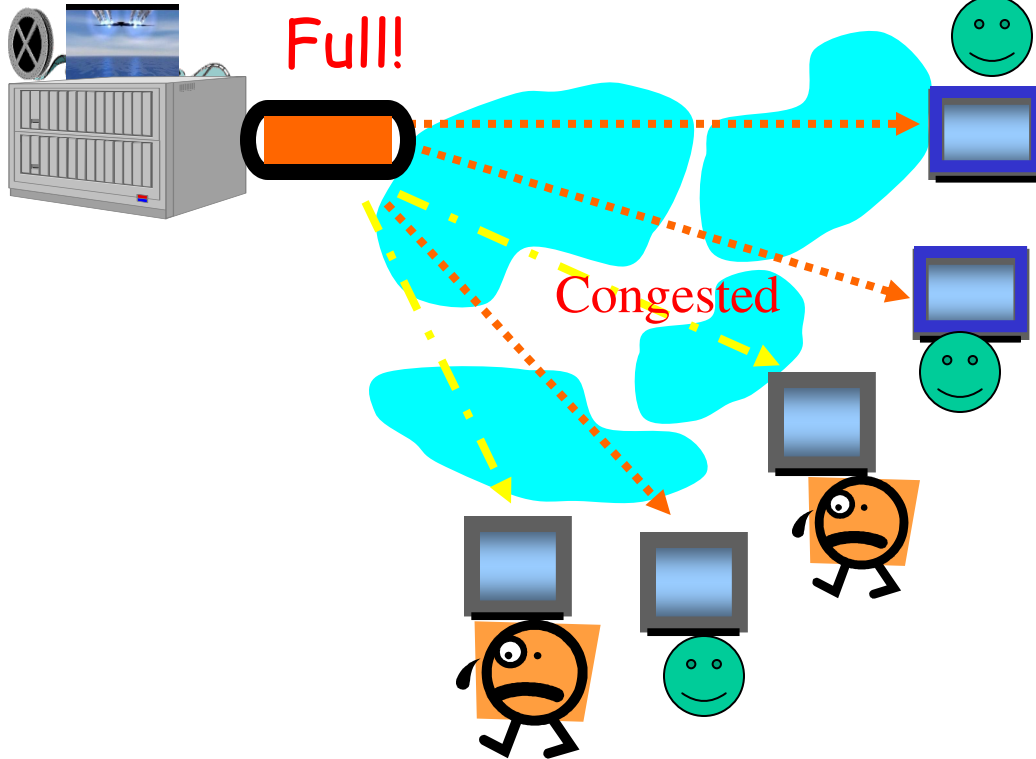


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# Outline

- ❑ Motivation
- ❑ Previous Work
- ❑ P2Cast: Peer-to-Peer Patching
- ❑ Overlay construction
- ❑ Providing continuous playback
- ❑ Performance evaluation
- ❑ Summary and future work

# Challenges of VoD Service



- ❑ High-bandwidth, long duration
- ❑ Large number of clients, asynchronous requests
- ❑ Client-server service model doesn't scale

# Previous Work on VoD Service

- ❑ IP multicast-based approach
  - Patching, Periodic Broadcast, Stream Merging
  - Problem: IP multicast has not been deployed widely
- ❑ Proxy Caching / CDN
  - Problems:
    - Infrastructure and maintenance cost
    - Still not scalable
- ❑ Peer-to-peer networks
  - Live streaming: SplitStream, PeerCast
  - P2P VoD: CoopNet, DirectStream

# P2Cast: Peer-to-Peer Patching Scheme

## ❑ Requirements:

- Unicast connection only

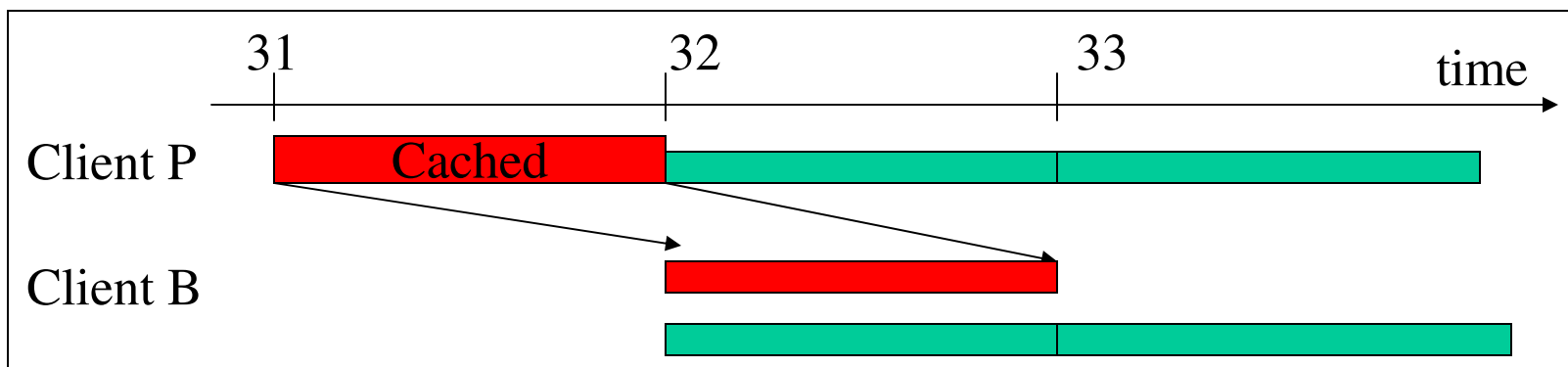
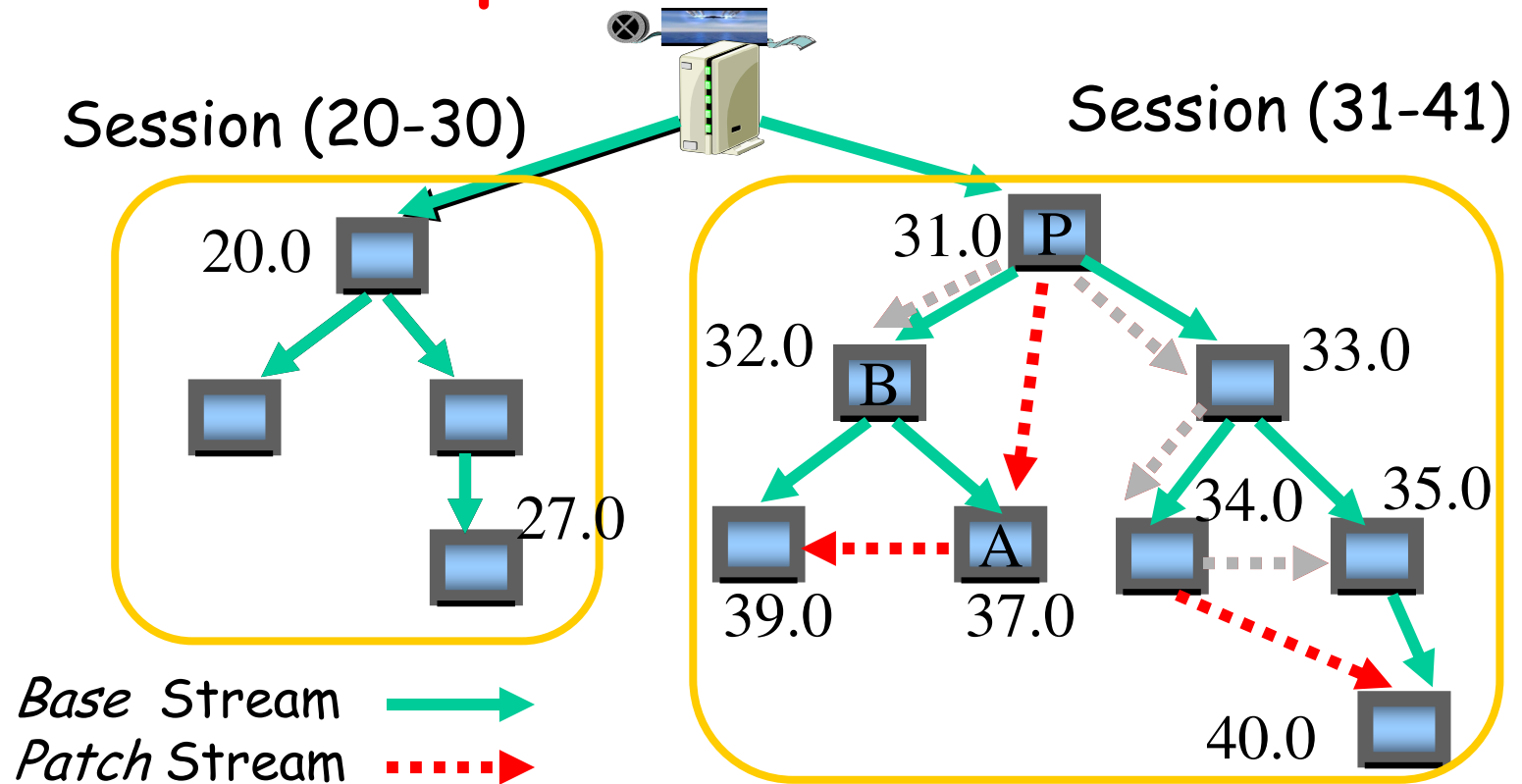
## ❑ Idea:

- Participating clients cooperate with each other
  - Forward stream to other clients
  - Cache and serve the initial part of the video to other clients

## ❑ Solution:

- Clients arriving close in time form a *session*
  - Closeness is determined by **Threshold**
- Establishing overlay among peers
  - Stream can be shared by all clients
- Missing *patch* can be retrieved from early-arrived clients

# A Snapshot of P2Cast

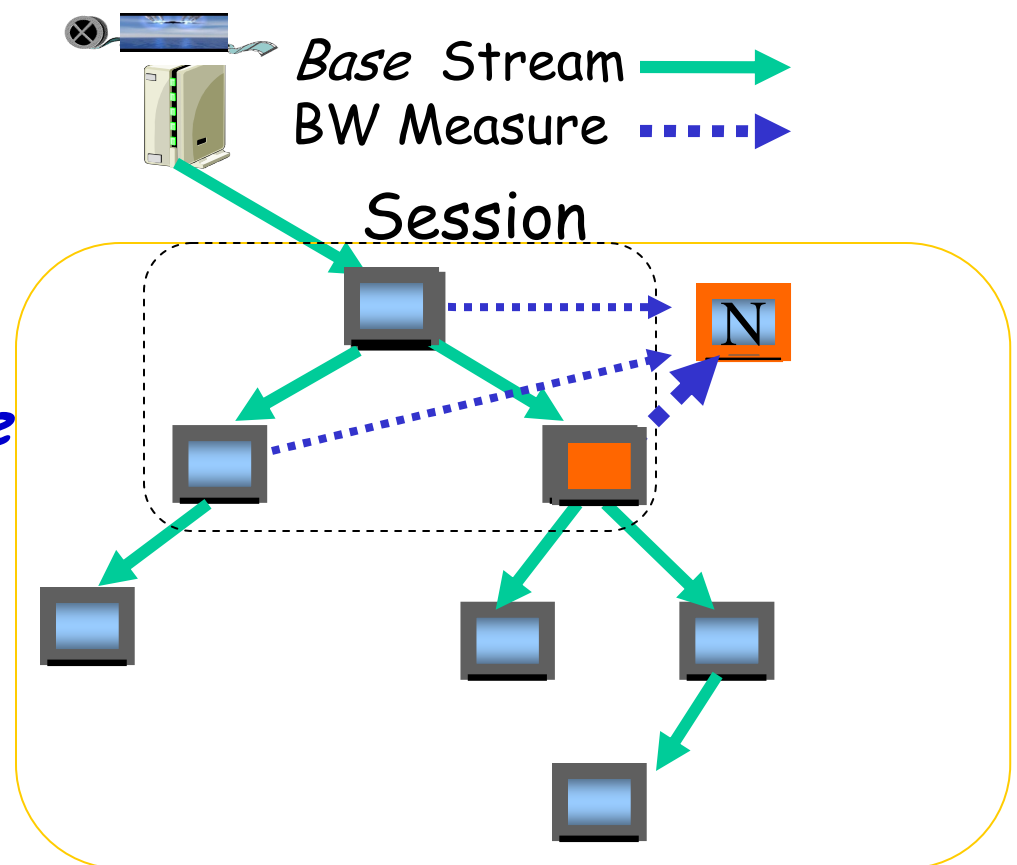


# Challenge for P2Cast

- ❑ Constructing application overlay appropriate for streaming
  - Supports sufficient Bandwidth to transmit the stream
  - Our solution: Best Fit (BF) Algorithm

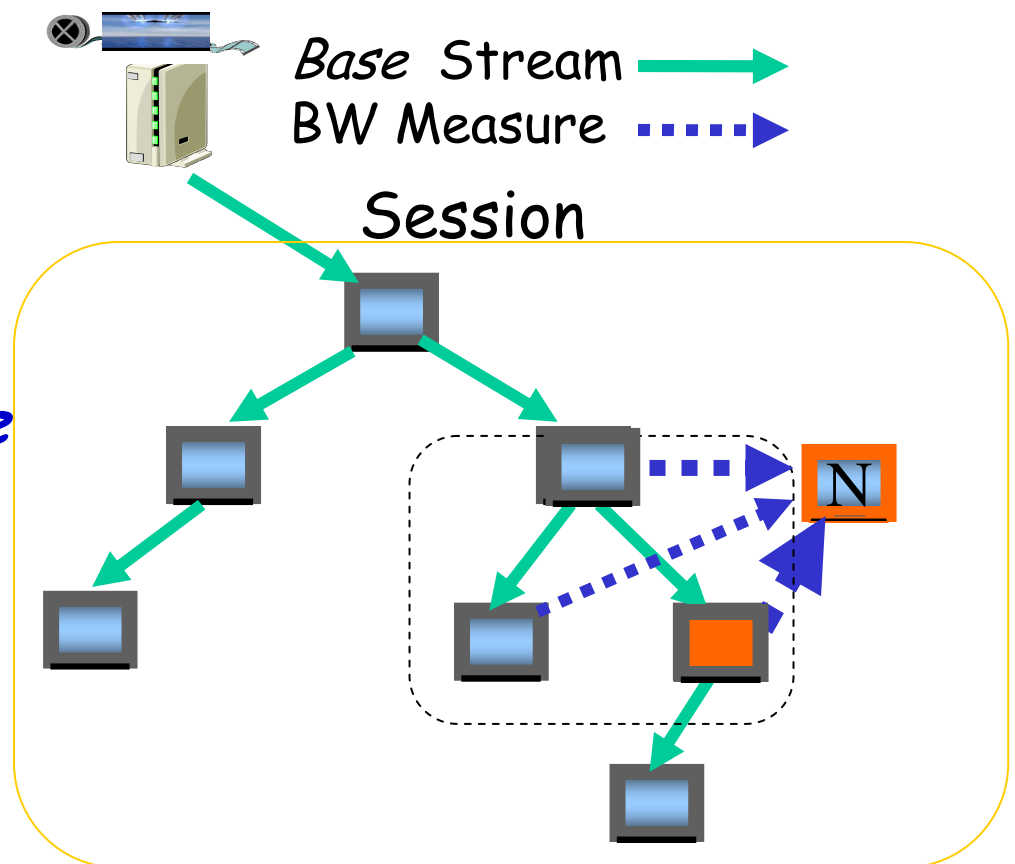
# "Best-fit (BF)" Overlay Construction

- Design Principles:
  - Bandwidth (BW)-first principle
  - Local information only principle
- Idea: Try to use **Maximum Available BW**
  - Measure BW in the unit of playback rate



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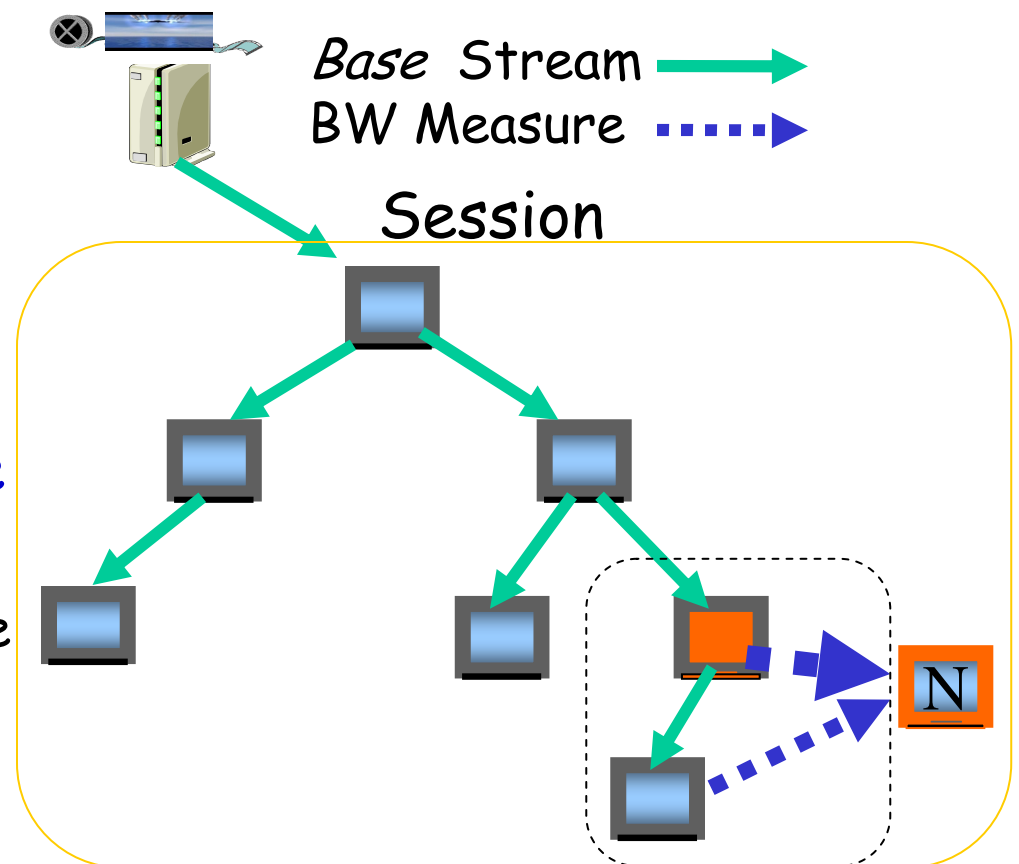
# "Best-fit (BF)" Overlay Construction

## □ Design Principles:

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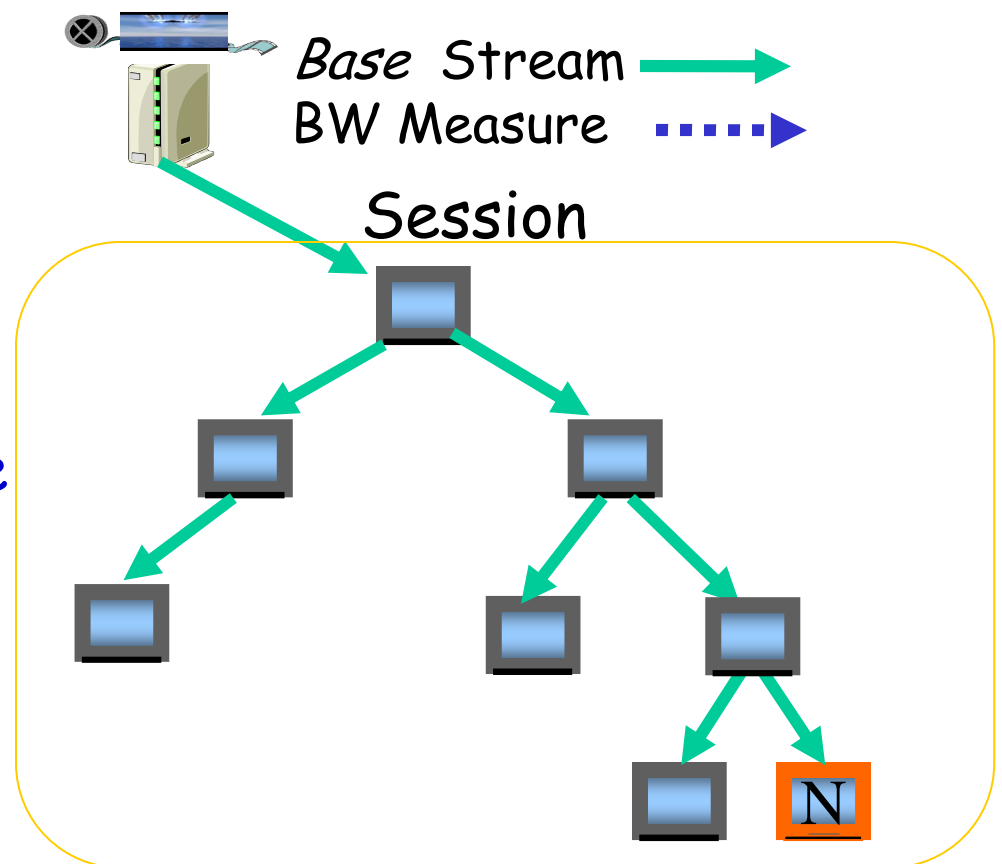
## □ Idea: Try to use Maximim Available BW

- Measure BW in the unit of playback rate



# "Best-fit (BF)" Overlay Construction

- Design Principles:
  - Bandwidth (BW)-first principle
  - Local information only principle
- Idea: Try to use Maximize Available BW
  - Measure discretized avail. bandwidth rounded to playback rate



# Two variations of BF: BF-Delay and BF-Delay-approx

## □ BF-delay

- Network delay information is used to break tie
  - Choose a closer client
  - Improve the *overall scalability*

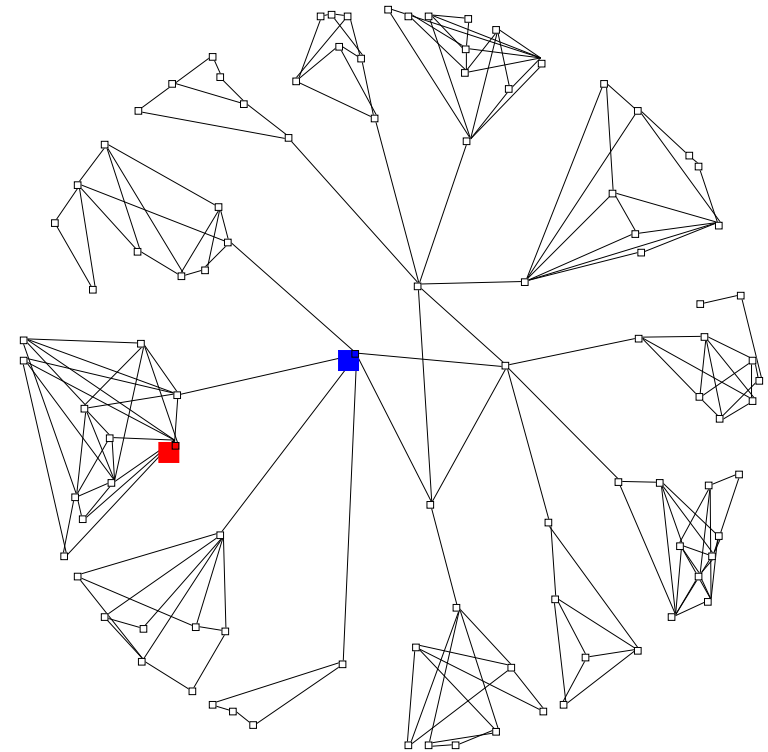
## □ BF-delay-approx

- Choose a peer with *sufficient* BW with delay info
  - Check Avail BW can sustain playback rate
  - Reduce measurement overhead
  - *Decrease joining delay*

# Performance Evaluation

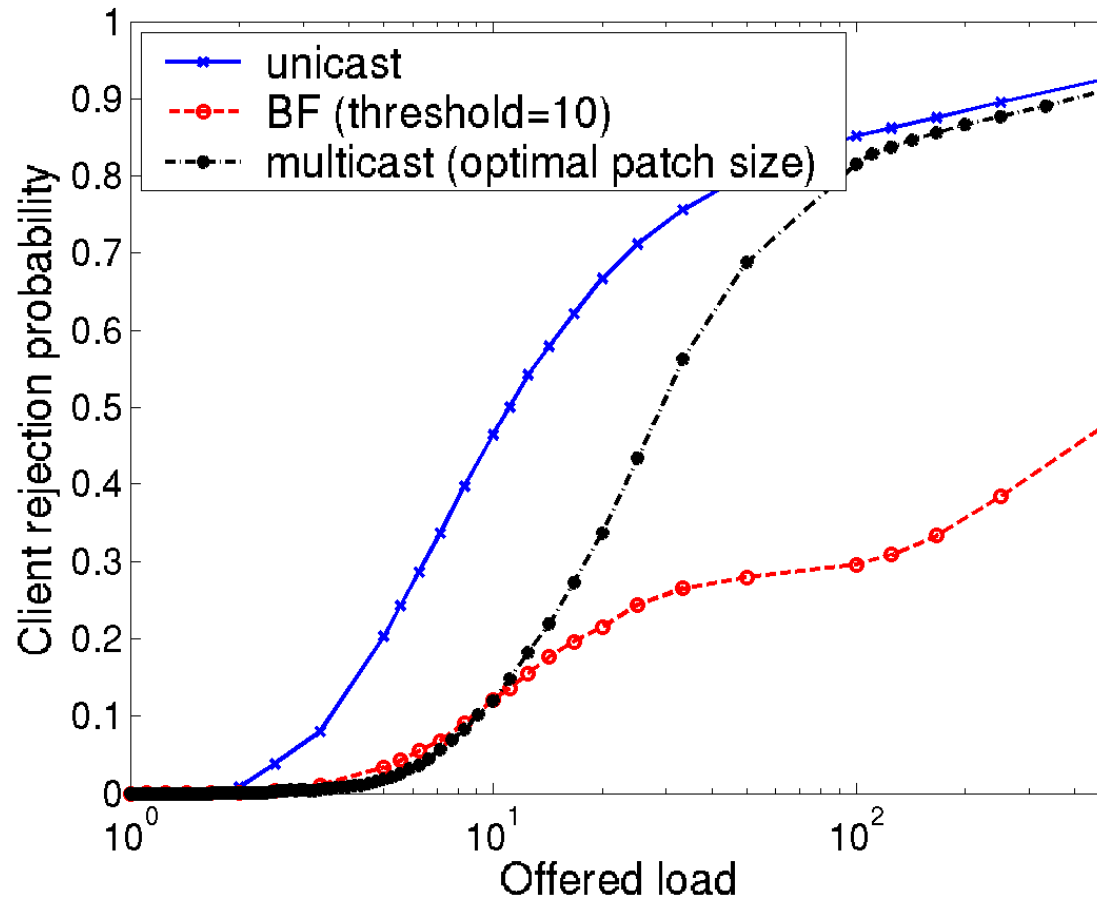
## □ Simulation settings:

- Transit-Stub network topology generated by GT-ITM (100 nodes)
- Node: Abstraction of a local network
- One Transit Network, 12 stub domains
- 95% confidence interval is within 5% of the point estimate



# Client Rejection Probability

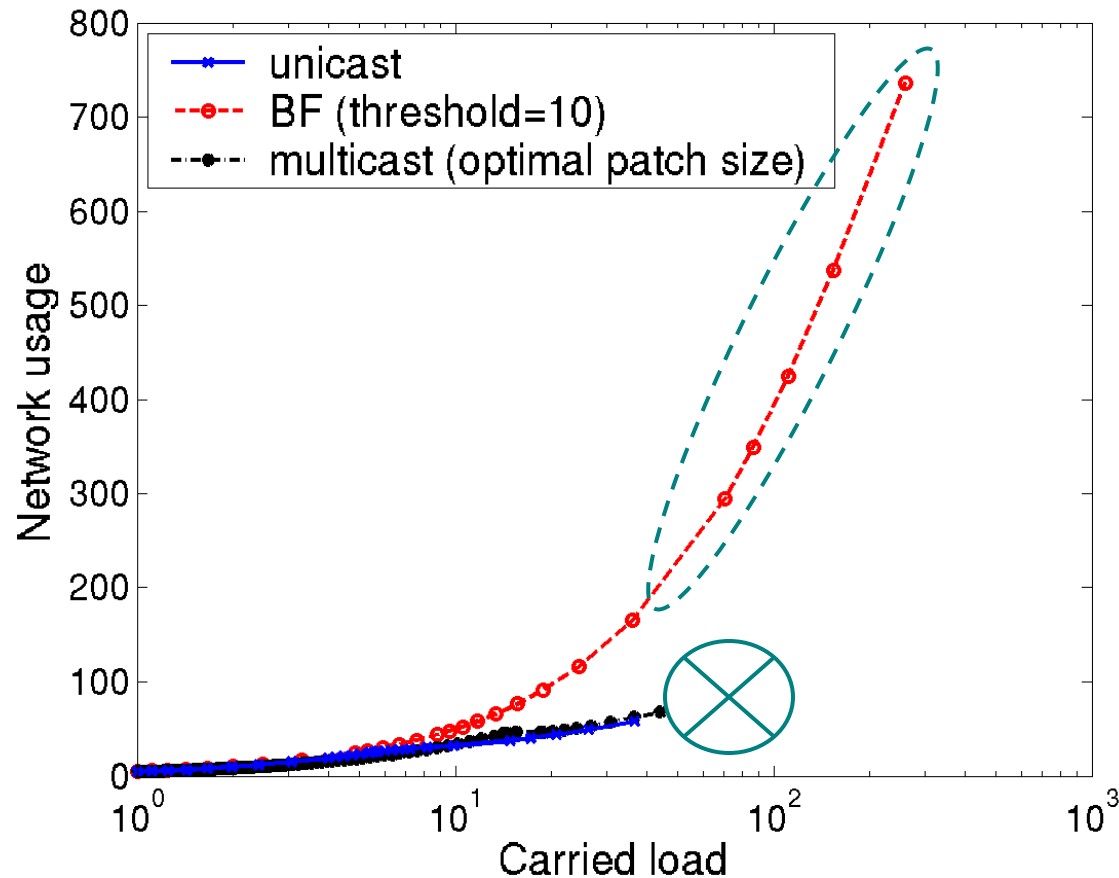
Server in Stub Domain



Rejection Probability in P2Cast using BF algorithm is smallest

# Network usage

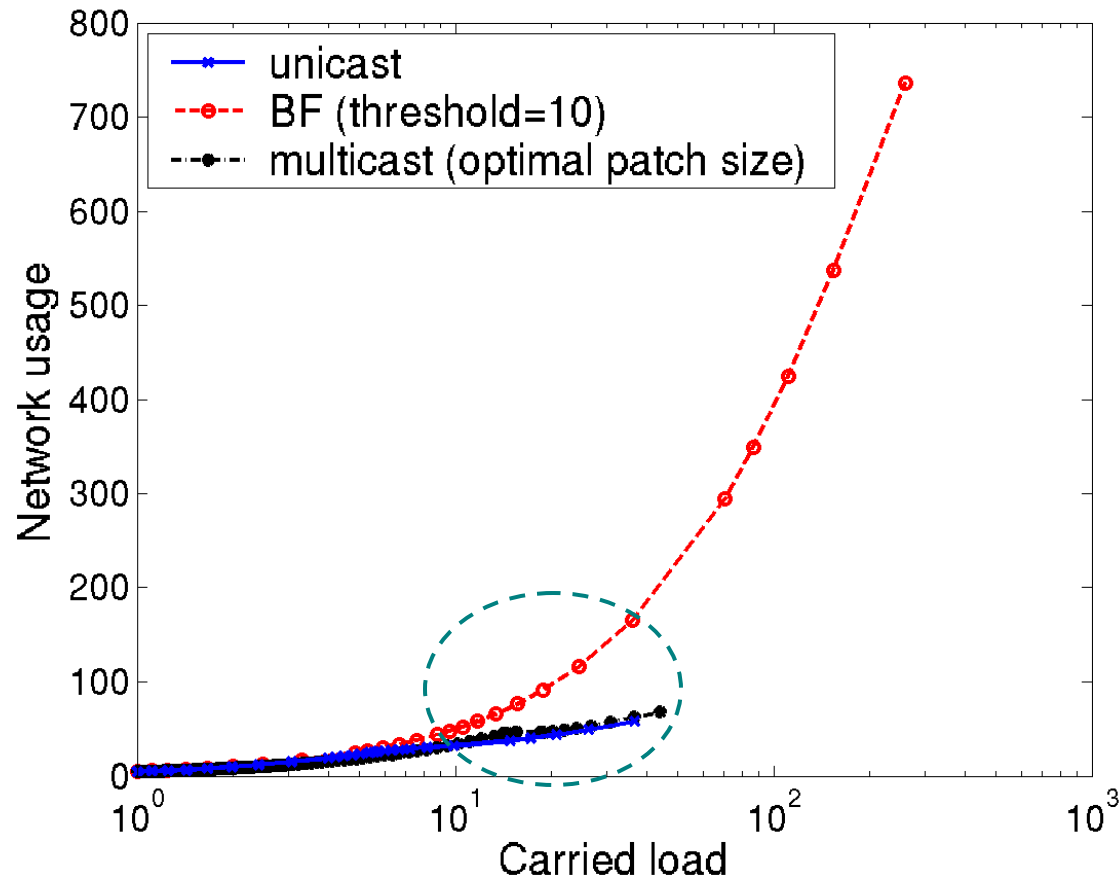
Server in Stub Domain



Network resources can be fully utilized in P2Cast with BF algorithm

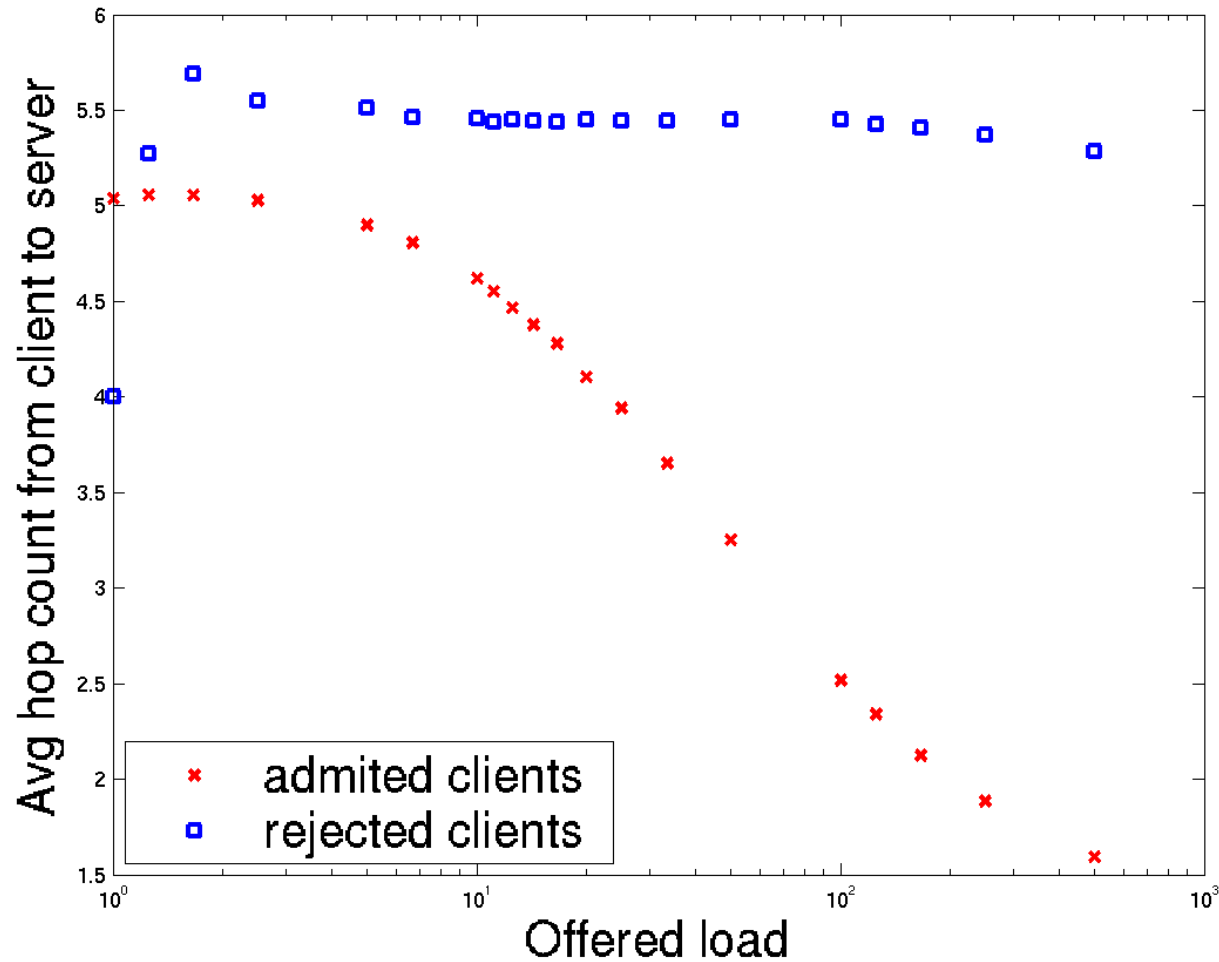
# Network usage

Server in Stub Domain

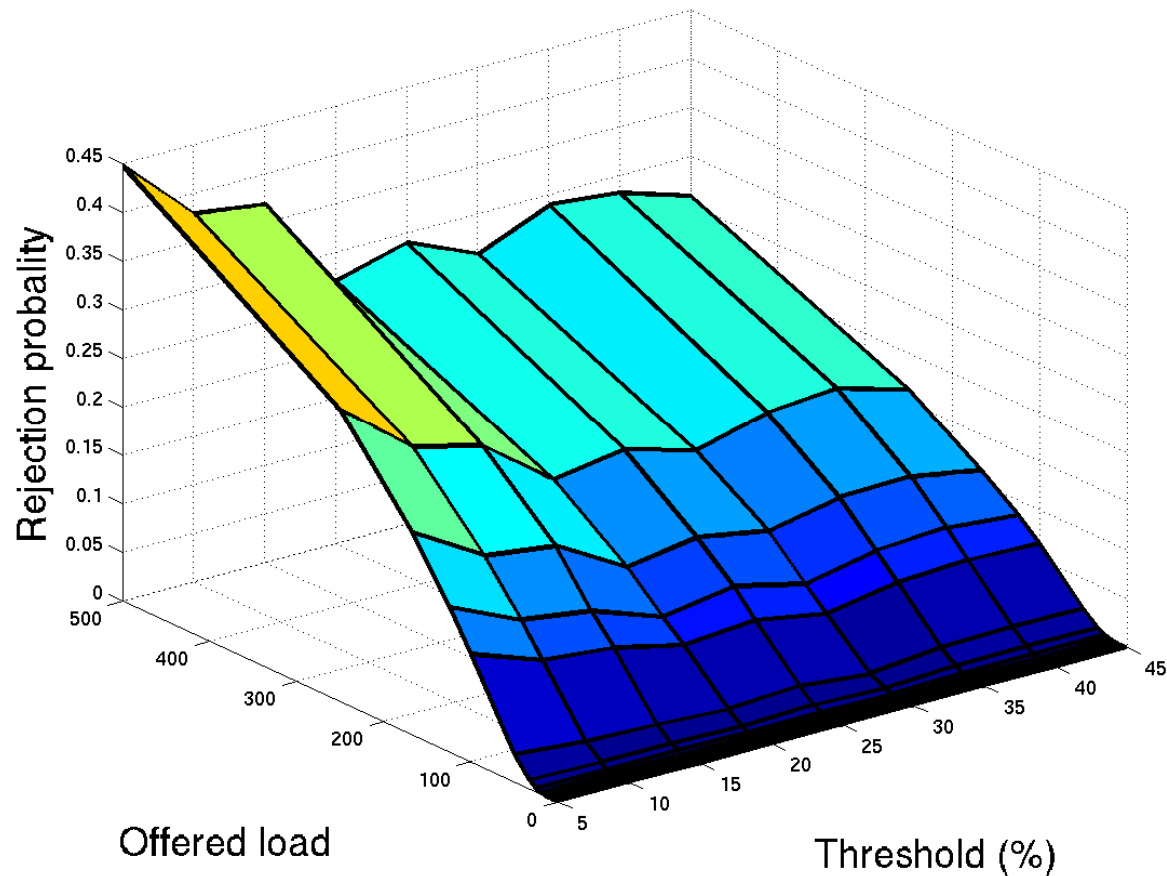


Why does P2Cast with BF algorithm have higher network usage than Unicast?

# Hop Count for Unicast



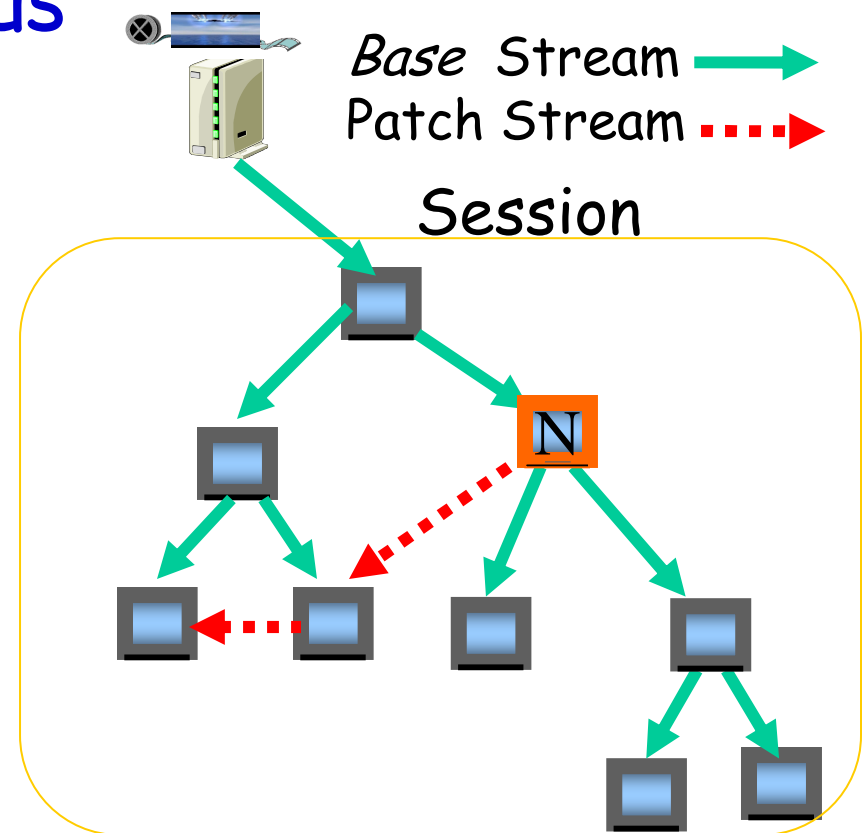
# Threshold impact on the performance of P2Cast



Larger Threshold → Better Scalability  
Small threshold is sufficient

# Challenge: providing continuous playback with departing clients

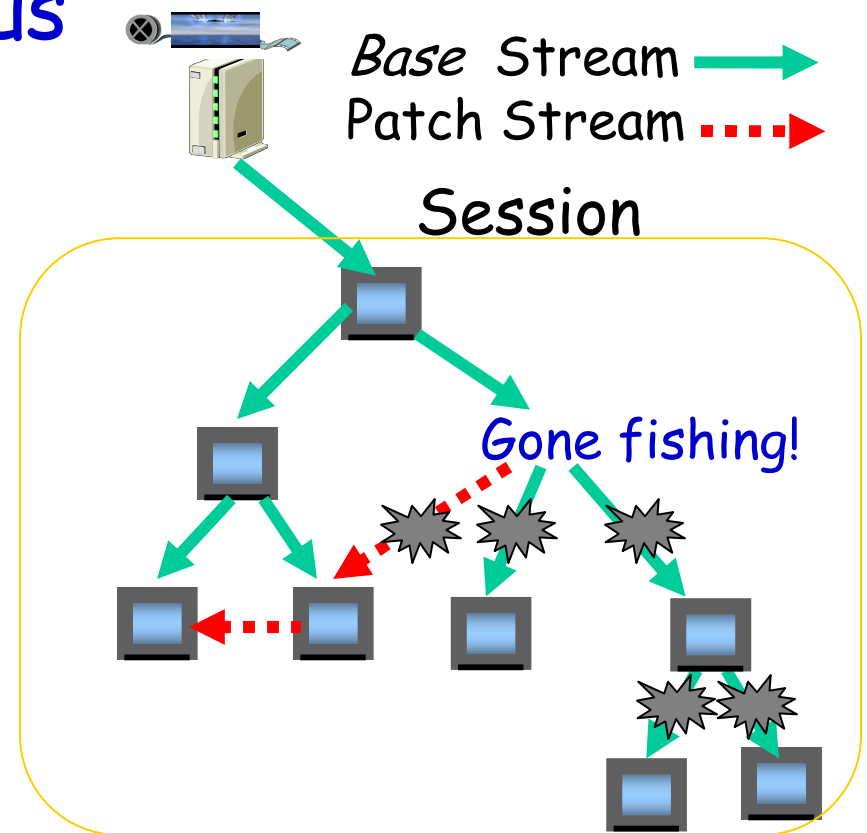
- Providing continuous stream playback without glitch
  - Resilient to disruption from early departing clients



# Challenge: providing continuous playback with departing clients

## □ Providing continuous stream playback without glitch

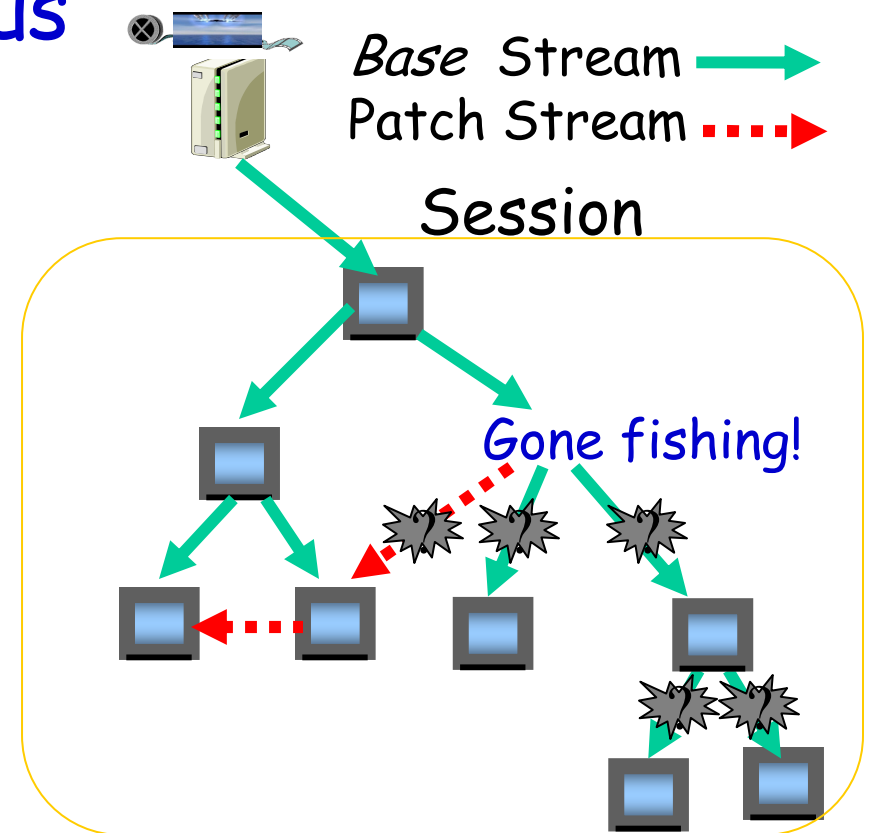
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# Challenge: providing continuous playback with departing clients

## □ Providing continuous stream playback without glitch

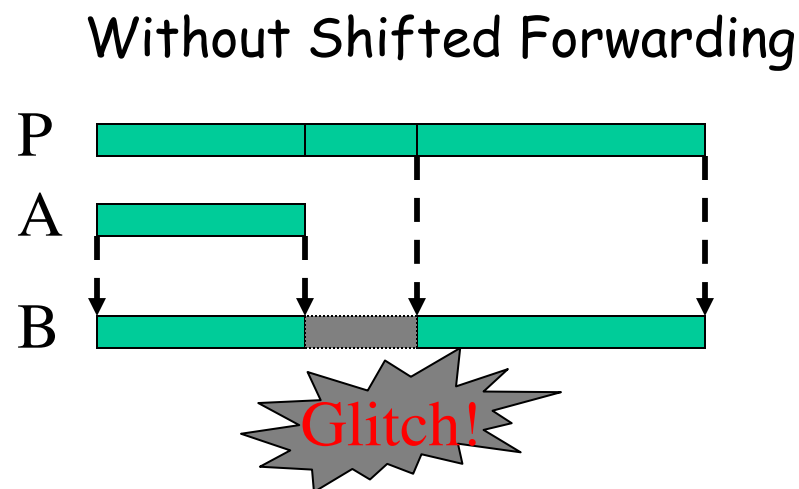
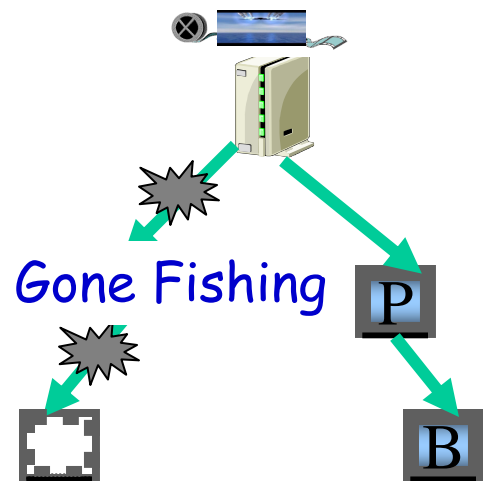
- Resilient to disruption from early departing clients



How Often? **0.004 for Patch**  
**0.0017 for Base**

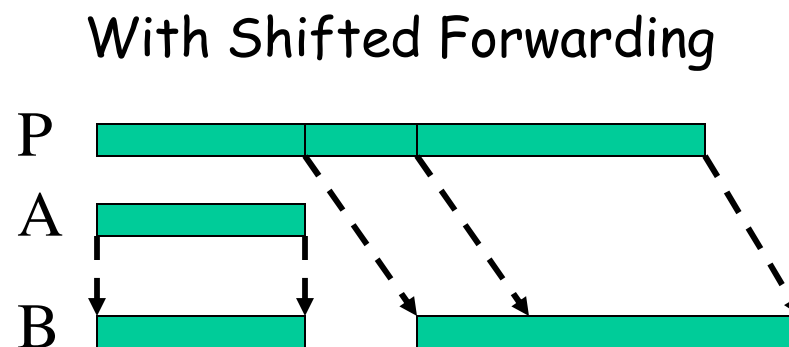
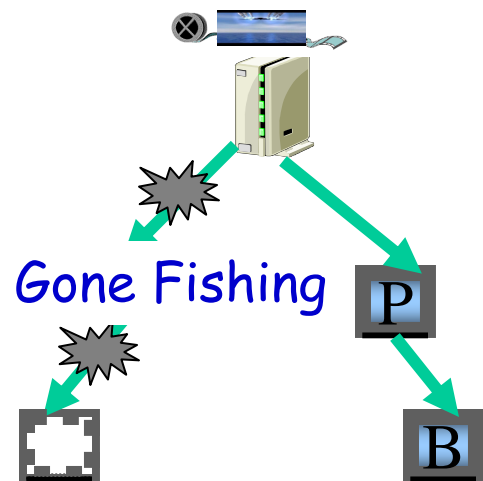
# Providing Continuous Playback

- Limit session size
- Buffering patch stream
- Shifted forwarding for base stream
  - Use of interval caching



# Providing Continuous Playback

- Limit session size
- Delay the start of playback for patch stream
- Shifted forwarding for base stream



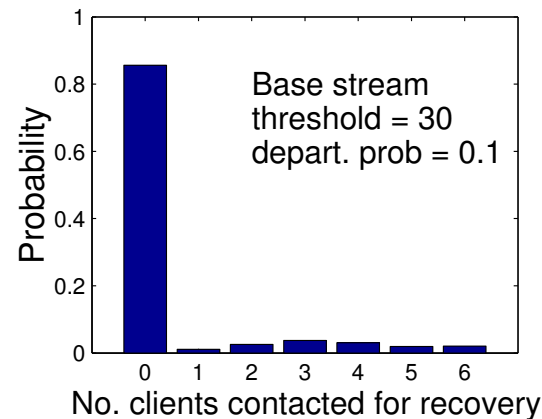
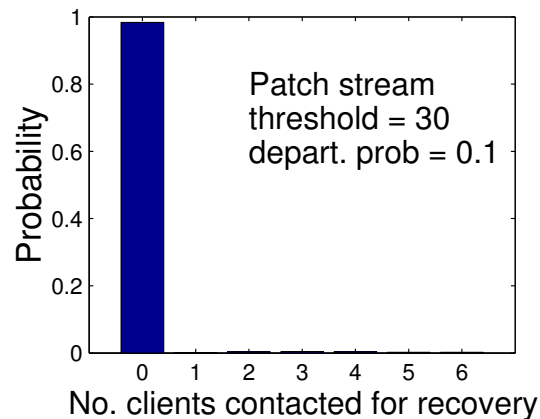
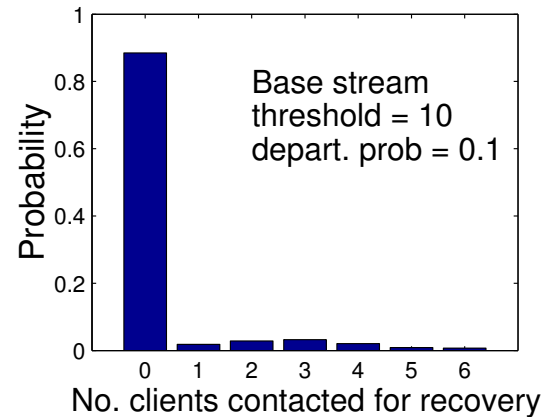
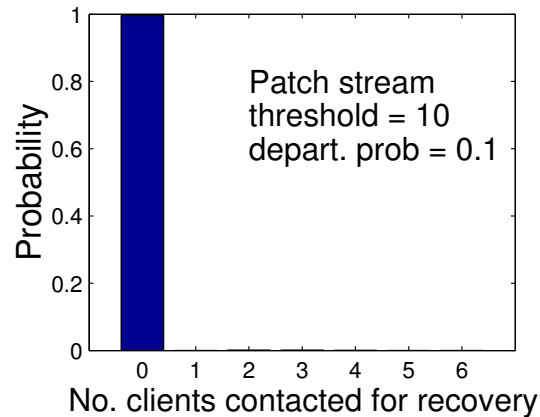
# Summary and Future work

- ❑ Proposed P2P patching scheme
  - Constructing application overlay appropriate for streaming
  - Providing continuous stream playback
- ❑ P2Cast outperforms:
  - unicast-based scheme and
  - native IP multicast-based patch scheme
- ❑ Experiment of P2Cast on the Internet

The End

# Backup Slides

# Probability of no. candidate clients contacted for recovery



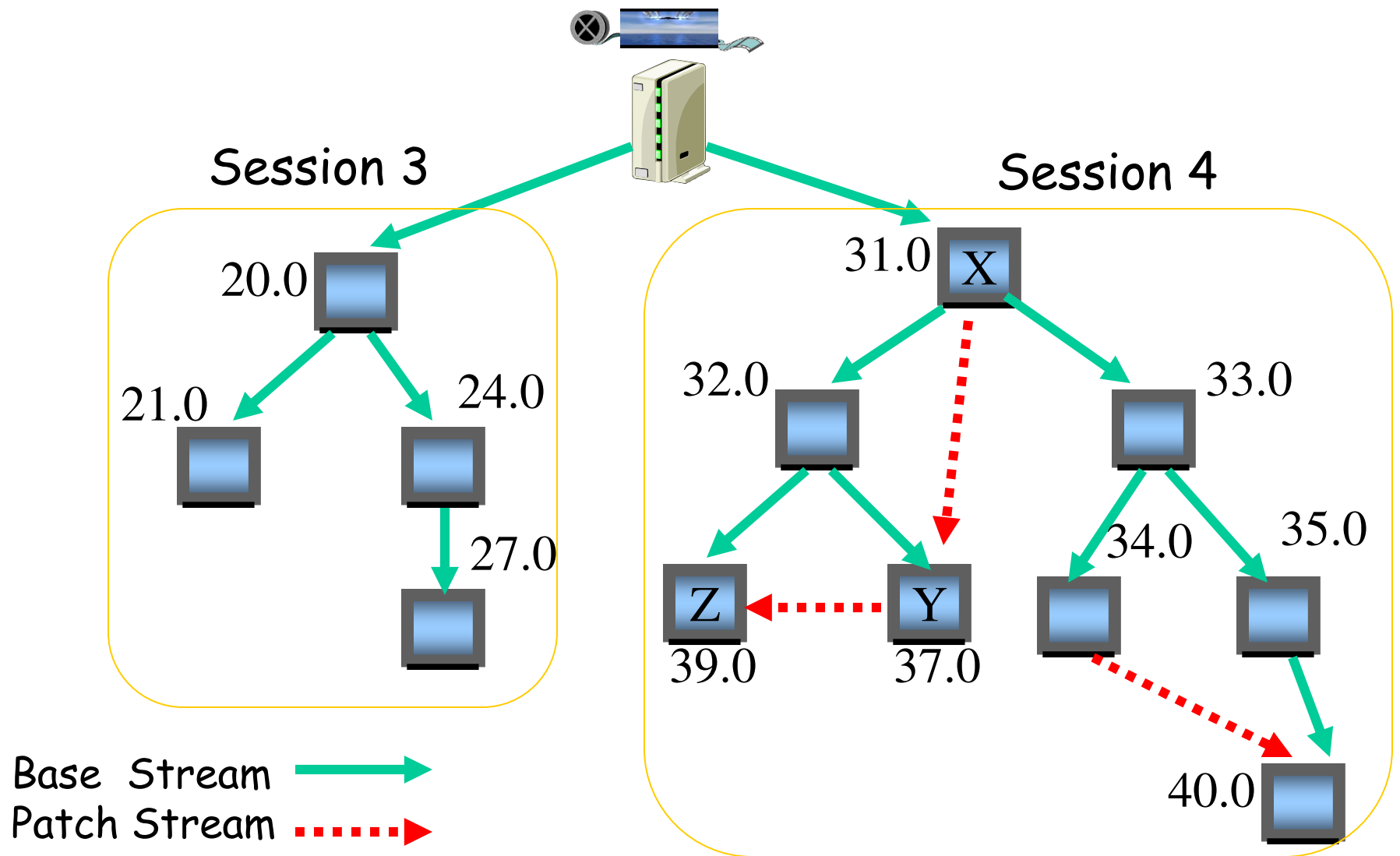
- Smaller Threshold shows better resilience to disruption
- Patch stream is relatively immune to disruption

# Expected Questions

1. Why not P2P PB or P2P streaming Merging?
2. Overhead of Avail-BW measurement?
3. Why not Narada + Patching Scheme?
4. Spatial Locality Information
  - Can be extended with Clustering mechanism given locality information



# A Snapshot of P2Cast at time 40



# Peers in P2Cast

## □ Role of Peers:

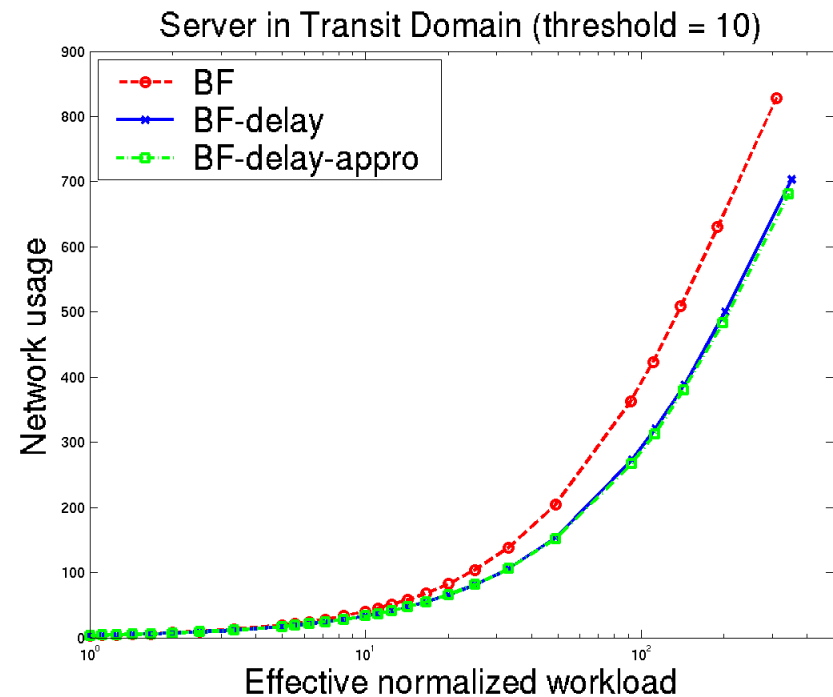
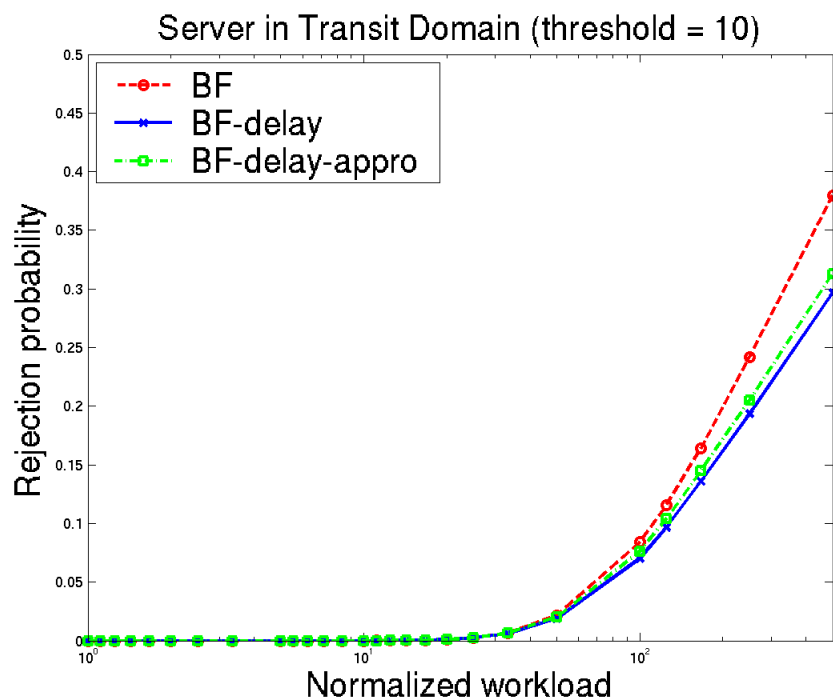
- Forward base stream to other clients
- Cache and serve the initial part of the video to other clients

## □ Benefit:

### Our Experimental Results show that

- P2Cast scales much better than traditional client-server unicast service
- If clients can cache > 10% of stream's initial part, it outperforms IP multicast-based patching scheme

# Comparison of BF, BF-delay, BF-Delay Approx



Exploiting "Delay" information make BF-delay and BF-delay approx perform better than BF

# "Best-fit (BF)" Overlay Construction algorithm

## □ Design Principles:

- Bandwidth-first principle
- Local information only principle

## □ Algorithm:

Step 1: Client N contacts a candidate parent P, starting with the server

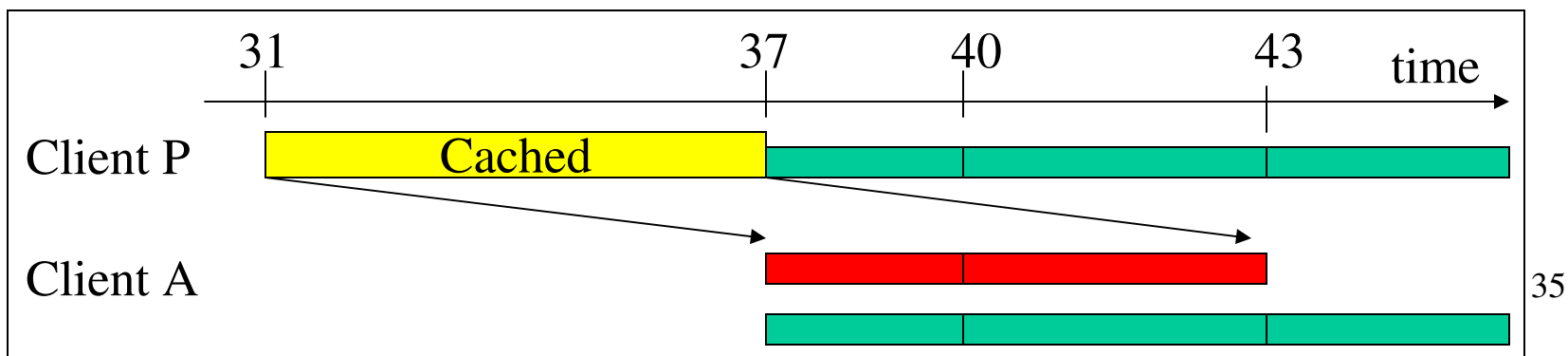
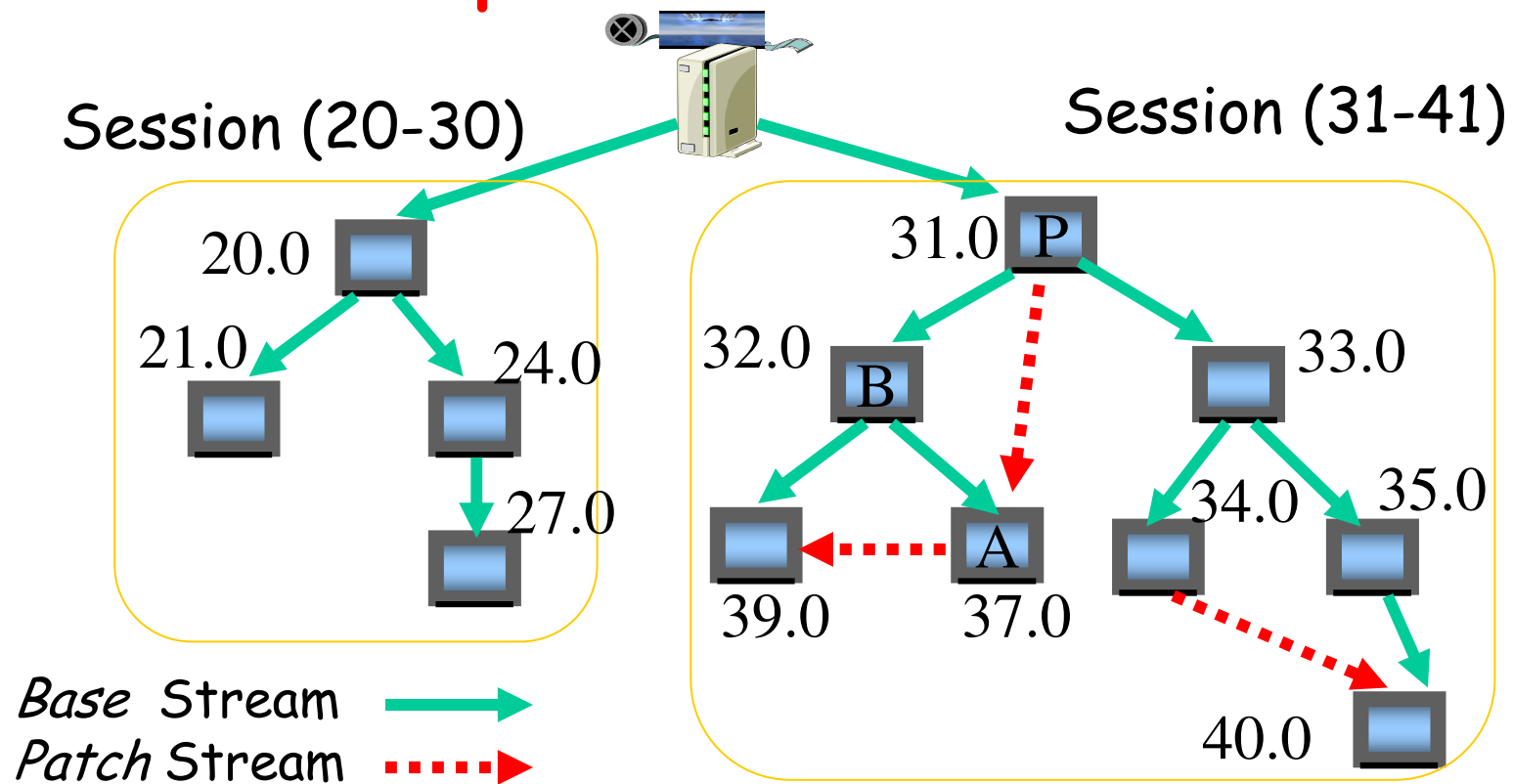
Step 2: Measure bandwidth from P to N,  $B(P,N)$  and  $B(C_i, N)$  where  $C_i \in \text{Children}(P)$

Step 3:  $C_{\max} = \operatorname{argmax}_{C \in \text{Children}(P)} B(C,N)$

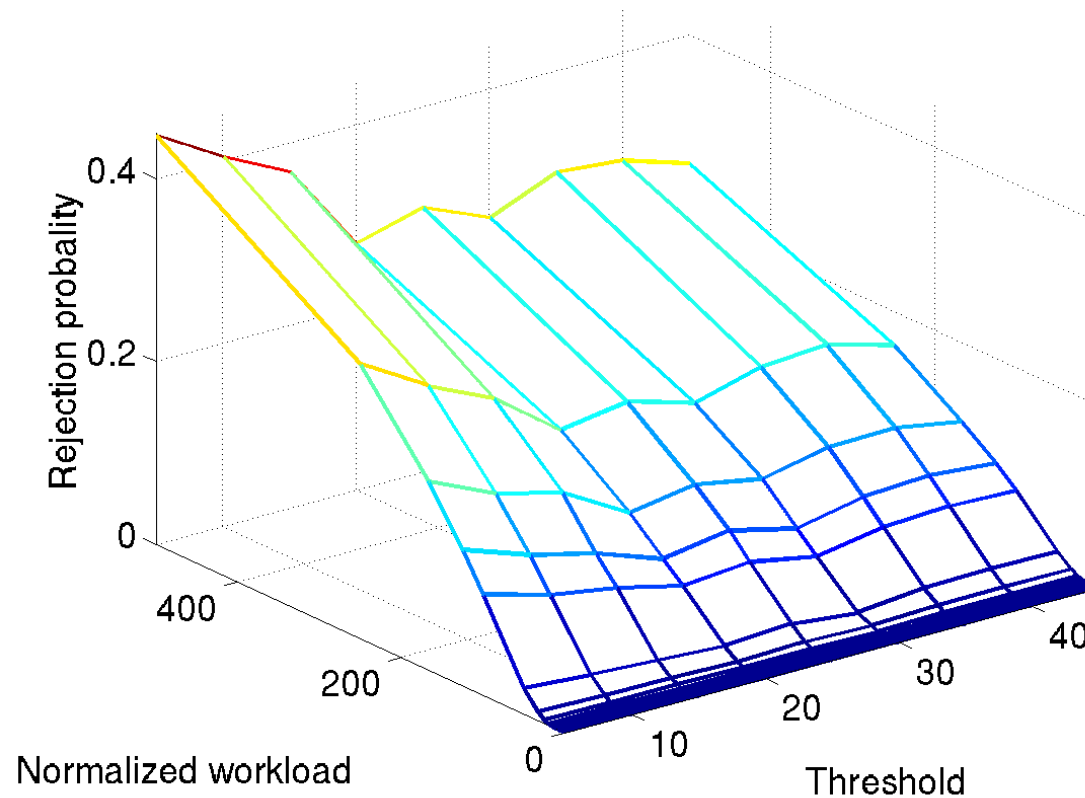
- If  $B(P,N) > B(C_{\max},N)$ , P provides service to N and Stop
- If  $B(P,N) \leq B(C_{\max},N)$ , N is re-directed to  $C_{\max}$  and starts from Step 1



# A Snapshot of P2Cast



# Threshold impact on the performance of P2Cast



Larger Threshold → Better Scalability  
Small threshold is enough

# Challenges for P2Cast

## 2. Providing continuous stream playback without glitch

- Resilient to disruption from early departing clients
- Our solution:
  - Limit session size
  - Delay the start of playback for patch stream
  - Shifted forwarding for base stream

